1. Suppose the set of atomic propositions is  $\{p_1, p_2\}$ . Consider the following transition system -



(Notation: For each state, the atoms written inside the curly braces next to the corresponding state are the atoms that are true at that state. For example,  $p_1$  is true at the top left state (initial state), whereas  $p_2$  is true at the bottom right state, and so on.)

Which of the following LTL formulas does this transition system satisfy?

- (a)  $Fp_2$
- (b)  $G(p_1 \vee p_2)$
- (c)  $(p_1Up_2) \vee G(\neg p_2)$
- (d)  $(p_1 \wedge p_2) \rightarrow X p_2$
- (e)  $G((p_1 \wedge p_2) \rightarrow Xp_2)$
- 2. Suppose p, q, r are three propositional atoms.
  - (a) Are the two formulas  $((p \ U \ q) \ U \ r)$  and  $(p \ U \ (q \ U \ r))$  equivalent? That is, whichever (infinite) word satisfies the first formula would satisfy the second and vice versa?
  - (b) Is  $(p \ U \ (q \lor r))$  equivalent to  $((p \ U \ q) \lor (p \ U \ r))$ ?
  - (c) Is  $((q \lor r) U p)$  equivalent to  $((q U p) \lor (r U p))$ ?
- 3. Suppose the set of atomic propositions is  $\{p_1, p_2\}$ . Consider the following NBA -



Show that the language of the above NBA is exactly the set of words satisfying the LTL formula  $F(\neg p_1) \wedge XGp_1$ .

- 4. Draw the NBA corresponding to LTL formulas  $p_1 U(\neg p_2), (\neg p_1) U p_2$ .
- 5. Let  $\phi, \psi$  and  $\chi$  be LTL formulas. We say two formulas are *equivalent*, written as  $\phi \equiv \psi$  if they define the same language. For each of the following, prove or disprove the equivalences:
  - (a)  $\boldsymbol{G}(\phi \wedge \psi) \equiv (\boldsymbol{G}\phi) \wedge (\boldsymbol{G}\psi)$
  - (b)  $GFG\phi \equiv FGF\phi$
  - (c)  $\mathbf{X}(\phi \mathbf{U}\psi) \equiv (\mathbf{X}\phi)\mathbf{U}(\mathbf{X}\psi)$
  - (d)  $(\phi U\psi)U\chi \equiv \phi U(\psi U\chi)$